Layering Protocol Verification: A Pragmatic Approach Using UVM

Rahul Chauhan (rchauhan@broadcom.com)
Gurpreet Kaire (gpsingh@broadcom.com)
Broadcom, Inc. San Diego, CA - USA
www.broadcom.com

Ravindra Ganti (rkganti@synopsys.com)
Subhranil Deb (sdeb@synopsys.com)

I. Introduction

Communication protocols are modeled as layers, and these layers are often labelled using the popular Open systems interconnection model. For transmission, the information flows from upper layers downstream to lower layers and for reception, the information flows upstream from lower

```
virtual task run_phase (uvm_phase );
  fork
    this.tx_driver ();
    this.rx_driver ();
  join_none
endtask : run_phase

virtual task tx_driver ();
  forever begin

seq_item_port.get_next_item (llc_frame )
    // -- Wait Before
    if (llc_frame.wait_before )

this.llc_frame_rcvd_ev.wait_ptrigger ();
    // -- Process frame for transmission
    if (!llc_frame.bypass_model )
      this.send_frame ();
    else
      this.send_corrupt_frame ();
    // -- Wait After
    if (llc_frame.wait_after )
      // -- Clear Existing Event
      if
(this.llc_frame_rcvd_ev. is_on())
        this.llc_frame_rcvd_ ev.reset ();

this.llc_frame_rcvd_ev.wait_ptrigger ();
    ...
    seq_item_port.done ();
  end
endtask : tx_driver

task send_frame (llc_frame );
  // -- Convert LLC to MAC
  this.convert_llc2mac (llc_frame );
  // -- Send to MAC passthru - sequence
  this.send _llc2mac (llc_frame );
  // -- Selective Enable Mechanism
  // -- Activate Timer for Flow Control
  if (!llc_frame.bypass_model )
    this.set_timer (llc_frame );
endtask : send_frame
```

```
        // Known Bad Tr,         No  Wait for
      response
          tr.expect_error = 1;
        if(msg_cnt == 2)
          ...
       `uvm_send(tr);
        get_response(rsp);
      end
    endtask : body

endclass : app_seq

class app_drv  extends uvm_component;
  task main_phase();
    forever bes3h3m49 0Zrg050
```


0 0 1 106.22 675.94 Tm 0 g 30 8.09 ET 0.275 rg 0.753 0n

```
//Use the monitor port to create wait
cond i tions for the sequences

//Wait task for waiting on one p                artic   u-
lar frame_kind from the monitor ports
task wait_for_frame(frame_kind_e
frame_kind);
   wait_for_frame_event(frame_kind);
    «
   process_frame_for_sequence();
endtask : wait_for_frame

//Implentation port write function i                    m-
plantation.

function write   _app();
    if (frame.frame_kind == frame_kind)
       «
      emit_frame_event();
endfunction : write_app

 function write_llc();
 endfunction : write_llc

function write_mac();
endfunction : write_mac

endclass : virtual sqr
```

Figure 6. ***Virtual Sequencer Class Code Block***

```
class virtual_seq      extends uvm_sequence;

 function new();
     get_handle_for_llc_passthru_seq();
     get_handle_for_mac_passthru_seq();
 endfunction : new

 task body();
    //Start application  sequence on a              p-
plication s      equencer

app_seq_1.start(p_sequencer.app_sqr);

app_seq_2.start(p_sequencer.app_sqr);

    //Wait for frame response for second
      application       seq by calling parent
      sequencer task
p_sequencer.
wait_for_frame(app_frame_kind);

    //  Optionally inject llc              error using
      pass  through       sequence in third
    application s      equence
    fork
       llc_passthru_seq.inject_error = 1;
    join_none
   app_seq_3.start(p_sequencer.app_sqr)

endtask : body

endclass : virtual_sequence
```

VI. Conclusion

The motivation for this paper is to share the concepts and simple techniques that were implemented and also share the benefits we achieved with the methodology. The techniques described in this paper can be extended to create even more robust and complex test pattern scenarios. The focus of this methodology was to have maximum controllability at every layer of abstraction while still having an automated test flow.

VII. References

[1] Synopsys, Inc., Beyond UVM: Creating Truly Reusable Protocol Layering.

[2] Accellera, Universal Verification Methodol-
R J \ 8 9 0 8 V H U ¶ V * X L G H